Game Search

CSCE 420 – Fall 2024

read: Ch. 5

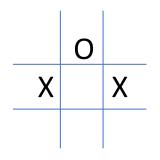
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Game Search

- games are useful to study for AI because they represent adversarial environments
 - the world state is not controlled solely by the agent
 - the world state can change because of actions by other agents (players)
 - different agents might have different objectives
 - this can lead to competitive behavior, or cooperative behavior
- there are many different kinds of games
 - simultaneous vs. sequential vs. iterated
 - single-player, two-player, multi-player
 - stochastic games with an element of chance
 - complete vs. incomplete information (*partially observable*)
 - also applies to economics: pricing of goods, auctions, contract negotiations...
- Of course, <u>DeepBlue</u> and <u>AlphaGo</u> are widely-recognized successes in Al, representing achievement of intelligent behaviour

Sequential Games

- multiple steps players take turns
- each player has a <u>utility function</u>
 - u_i(s) (where *i* is the player, and *s* is a game state)
 - +1 for win; -1 for lose; 0 for draw (tic-tac-toe); 0 for non-terminal states
 - money (poker)
 - rewards for achieving goals cost of actions or resources used
- simplest form: <u>2-player, 0-sum games</u>
 - $\Sigma_i u_i(s) = 0$ or $u_1(s) = -u_2(s)$
- examples: tic-tac-toe, checkers, chess...

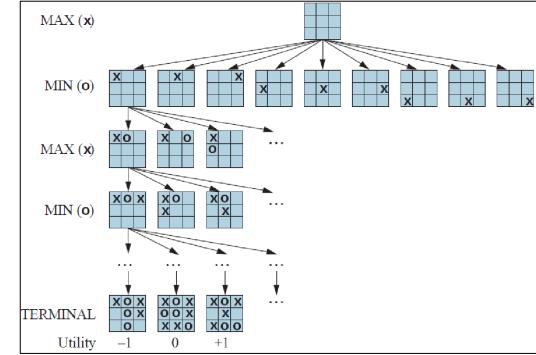


- in a 2-player, 0-sum game like tic-tac-toe, how can we decide what move to make?
- method 1: write a bunch of rules that encode a *strategy*
- method 2: use systematic search
 - use *look-ahead* for each possible action to imagine what opponent response might be
 - key idea: we can anticipate what move the opponent will make, because their utility is assumed to be the opposite of ours
 - thus the opponent will change the game in the way that is best for them, which is worst for us
 - *recursion*: of course, to simulate the opponent's reasoning, they will have to consider our response to their response, and so on...

- recall that u_i(s)=0 for non-terminal states
- label alternating levels in search tree as <u>max nodes</u> and <u>min nodes</u>
- define *minimax* value for each state s as follows:

 $u_i(s) \text{ if } s \text{ is a terminal state}$ minimax(s) = - max { minimax(s') for s' \in succ(s) } if s is a max node min { minimax(s') for s' \in succ(s) } if s is a min node

- decision at root node: argmax { minimax(s') for s'∈ succ(s) }
 - i.e. choose the action that leads to the successor with highest score, which has the highest expected payoff

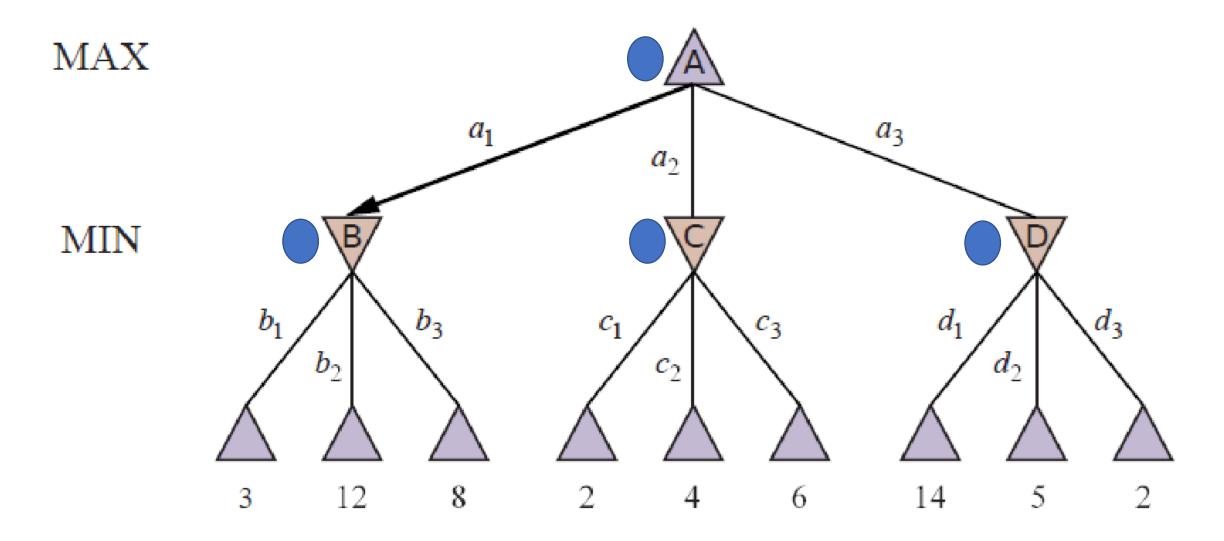


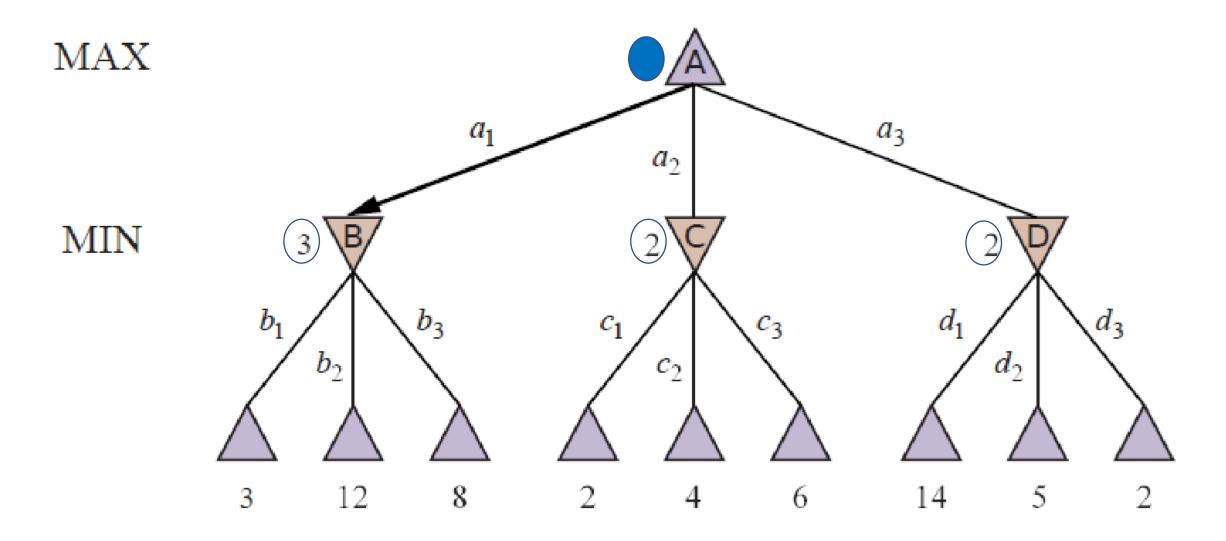
 $\begin{array}{l} \textbf{function } \texttt{MINIMAX-SEARCH}(game, state) \textbf{ returns } an \ action \\ \texttt{player} \leftarrow game.\texttt{TO-MOVE}(state) \\ \textit{value, } move \leftarrow \texttt{MAX-VALUE}(game, state) \\ \textbf{return } move \end{array}$

function MAX-VALUE(game, state) returns a (utility, move) pair if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null $v \leftarrow -\infty$ for each a in game.ACTIONS(state) do $v2, a2 \leftarrow MIN-VALUE(game, game.RESULT(state, a))$ if v2 > v then $v, move \leftarrow v2, a$

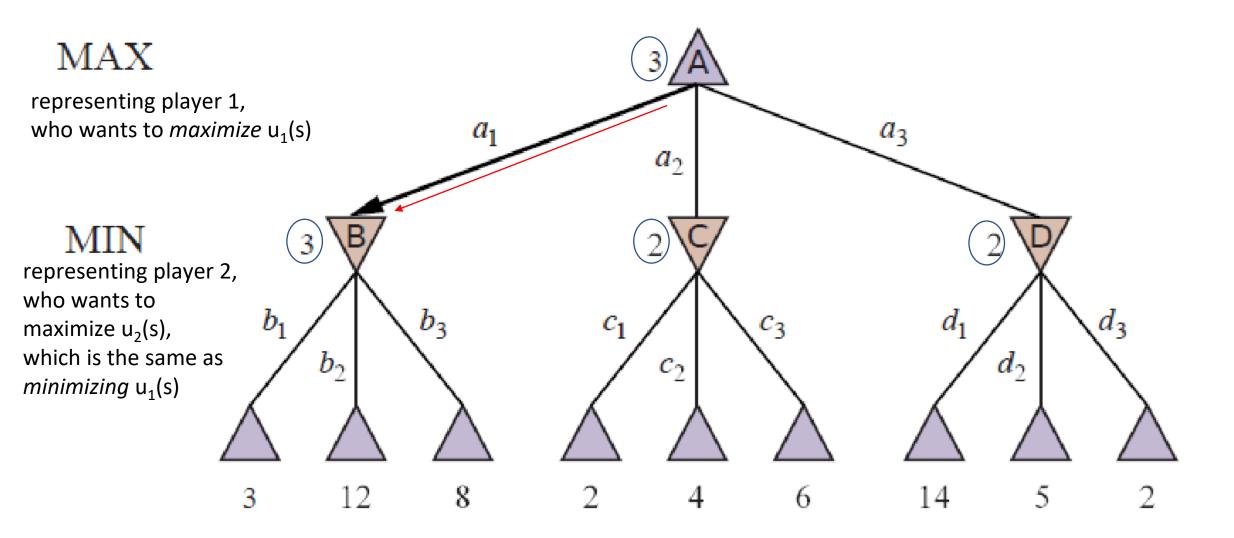
return v, move

function MIN-VALUE(game, state) returns a (utility, move) pair if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null $v \leftarrow +\infty$ for each a in game.ACTIONS(state) do $v2, a2 \leftarrow MAX-VALUE(game, game.RESULT(state, a))$ if v2 < v then $v, move \leftarrow v2, a$ return v, move *double-recursion: each function calls the other*





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- note: this only determines next move (by player 1)
- then player 2 chooses an action
- then we have to recompute the game tree from that state to decide the next move
- minimax does not determine the entire sequence of play; you cannot force the choices of the other player
- we *assume* the opponent will make optimal choices (for them)
- what happens if they make a sub-optimal move (e.g. a mistake)?

Complexity of Game Search

- the problem with applying Minimax to most games is that the search space is too large
 - estimates for chess: avg game=70 moves, avg branching factor=35, state space = ~35⁷⁰ = ~10¹⁰⁸
 - so we can't search all the way to leaves (end-games) where utility is defined to propagate the minimax values back up
- solution 1: use intelligent *pruning* to reduce the search space
 - sometimes we can infer parts of the space that do not need to be searched

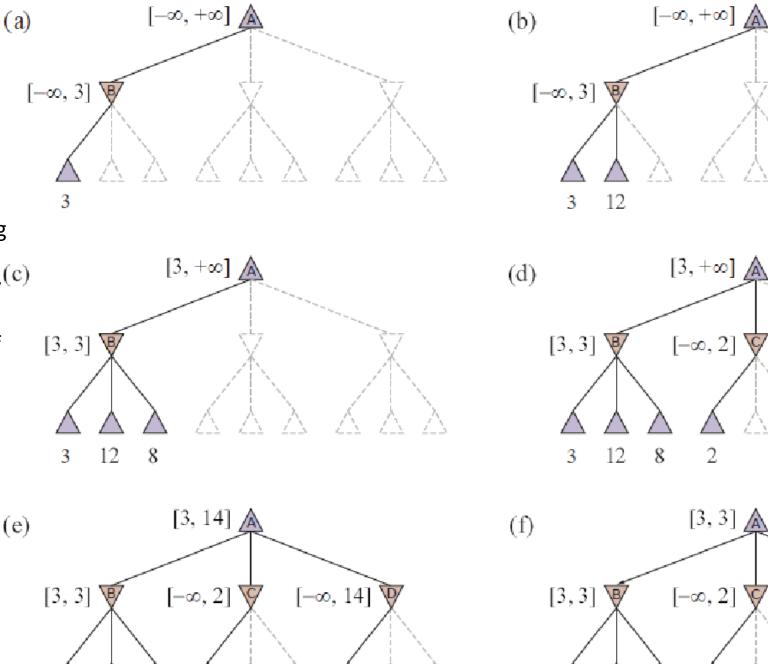
α/β -pruning

- at each node, keep track of 2 additional values α , β (along with minimax value)
 - α is the best possible value for any max node above so far (initially - ∞)
 - β is the best possible value for any min node above so far (initially + ∞)
- as we process children, update these params
 - at max nodes, update α : α =max{ α , minimax(s')} for each s' \in children(s)}
 - at min nodes, update β : β =min{ β ,minimax(s')} for each s' \in children(s)}
- pruning condition:
 - at min nodes: when v< α (i.e. best choice of parent max node)
 - at max nodes: when $v>\beta$ (i.e. best choice of parent min node)
- equivalently: when interval of v at node no longer overlaps interval of parent 2/19/2023

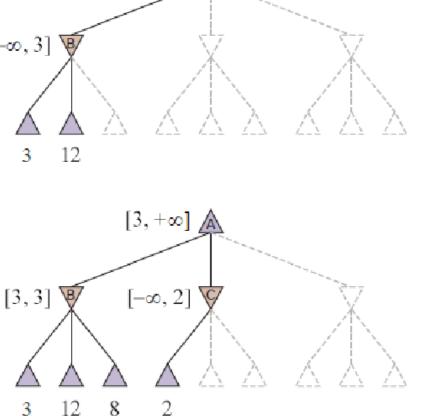
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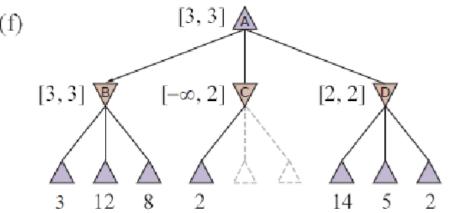
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(this example is for a simplified version of the alpha-beta pruning algorithm where we initialize minimax value v $^{(c)}$ to the range $[-\infty,\infty]$ at every node (instead of passing α and β in as parameters), and the pruning condition is evaluated by checking the overlap between the range of each node and (e) it's parent)



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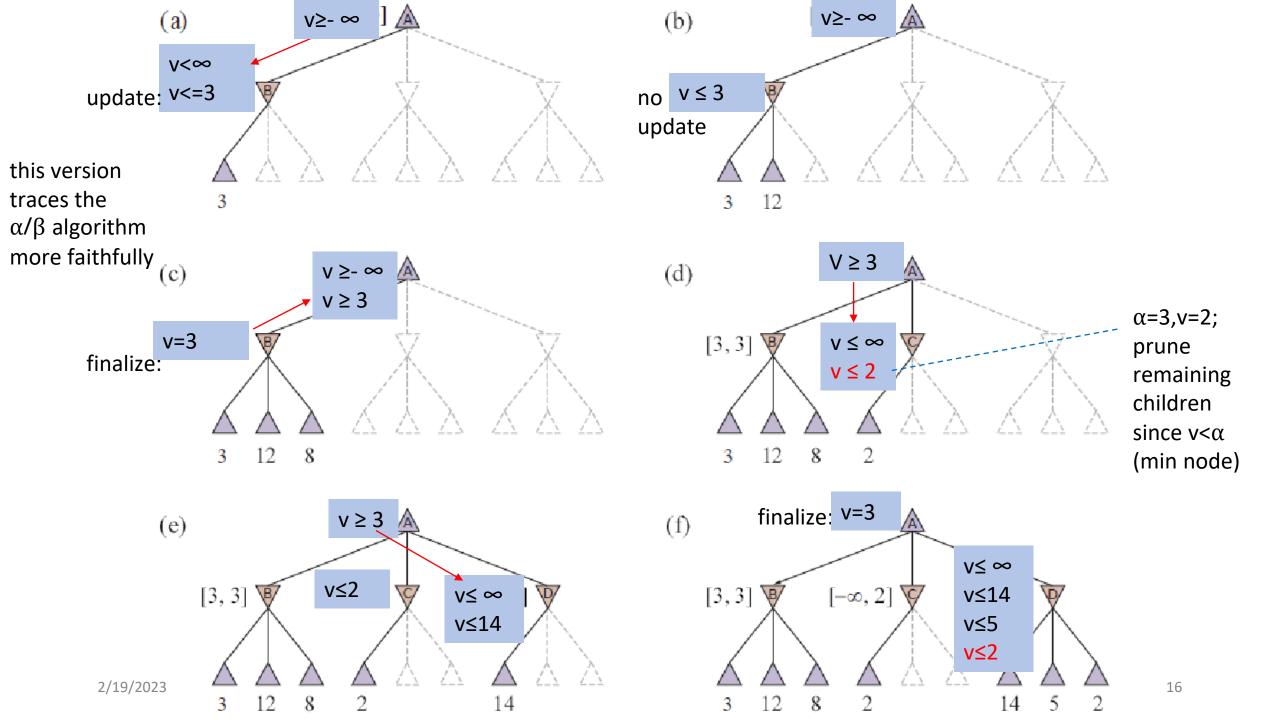
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function ALPHA-BETA-SEARCH(game, state) returns an action
                                    player \leftarrow game.To-MOVE(state)
                                    value, move \leftarrow MAX-VALUE(game, state, -\infty, +\infty)
                                    return move
                                function MAX-VALUE(game, state, \alpha, \beta) returns a (utility, move) pair
                                    if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null
                                    v \leftarrow -\infty
                                    for each a in game.ACTIONS(state) do
                                       v2, a2 \leftarrow \text{MIN-VALUE}(game, game.\text{RESULT}(state, a), \alpha, \beta)
                                       if v2 > v then
max nodes update \alpha \longrightarrow \begin{bmatrix} v, move \leftarrow v2, a \\ \alpha \leftarrow MAX(\alpha, v) \end{bmatrix}
if v \ge \beta then return v, move \leftarrow f prune if score becomes greater than upper-bound of parent's interval, since parent would never
                                    return v, move
                                                                                           choose this branch
                                function MIN-VALUE(game, state, \alpha, \beta) returns a (utility, move) pair
                                    if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null
                                    v \leftarrow +\infty
                                    for each a in game.ACTIONS(state) do
                                        v2, a2 \leftarrow \text{MAX-VALUE}(game, game.\text{RESULT}(state, a), \alpha, \beta)
                                       if v2 < v then
\begin{array}{c} \text{min nodes update } \beta \longrightarrow \begin{array}{c} v, \ move \leftarrow v2, \ a \\ \beta \leftarrow \text{MIN}(\beta, v) \end{array} \\ \hline \mathbf{if} \ v \leq \alpha \text{ then return } v, \ move \end{array}
                                    return v, move
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Complexity of Game Search

- solution 2: use a depth-limit while searching a game tree
 - need a *board-evaluation function* to assign scores to internal nodes (or non-terminal states, or non-end-games)
 - the value estimates the probability of winning or expected payoff from each state (heuristically)
 - the computer can then perform Minimax (possibly with α/β -pruning) down to a fixed level, apply the board evaluation function, and propagate values upward
 - choose depth limit based on time available (and CPU speed)
 - expressed as number of "ply" (moves, or levels)
 - 2-6 ply (a few sec): rudimentary chess performance (amateur skill level)
 - 6-10 ply (a few min): much better moves due to deeper search/look-ahead

Board Evaluation Functions

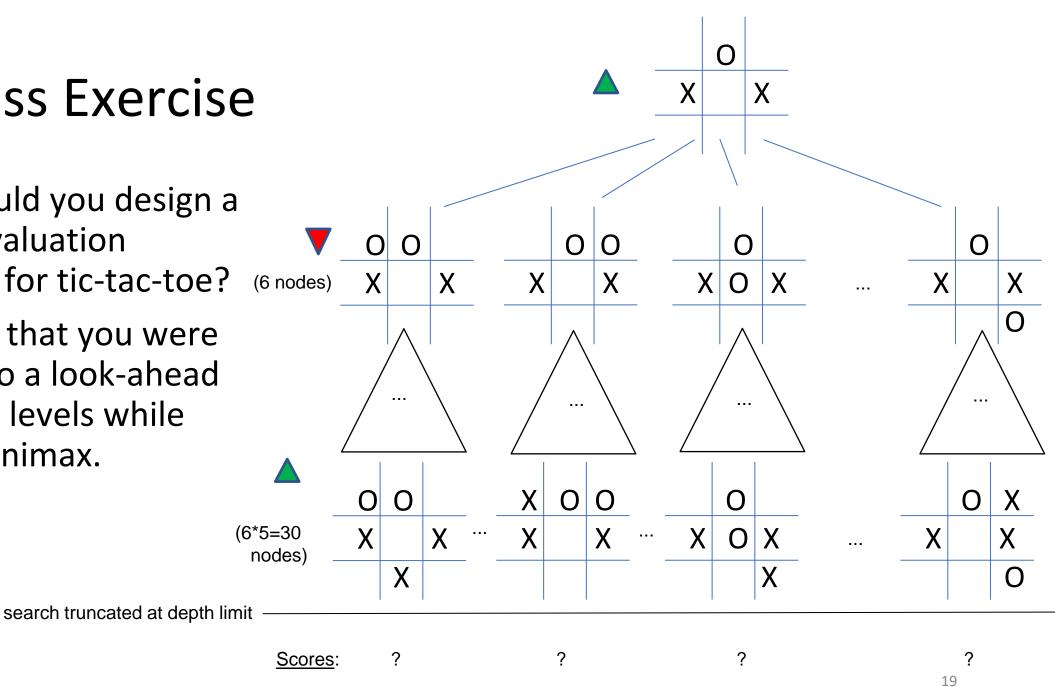
- a board evaluation function must guess the value (probable outcome) of each state
- they are typically based on *features*
- examples from chess:
 - piece differential (#PlayerPieces #OpponentPieces)
 - material value (pawn=1, knight/bishop=3, rook=5, queen=9)
 - center control
 - # of pieces threatened or constrained
 - patterns or special arrangements of pieces

 $Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$



In-class Exercise

- How would you design a • board evaluation function for tic-tac-toe? (6 nodes)
- Suppose that you were limited to a look-ahead of only 2 levels while doing minimax.



Board Evaluation Functions

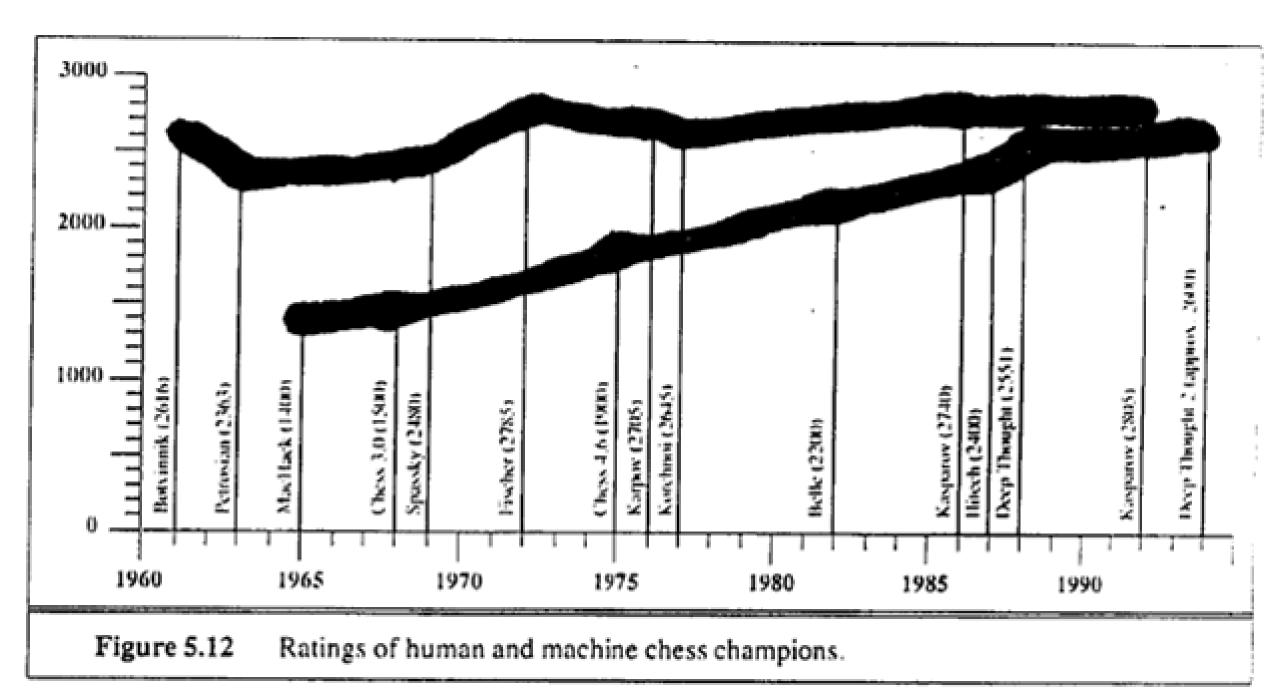
- problems with using board evaluation functions
 - non-quiescence
 - board evaluation function should only be applied to quiescent states, where the value has stopped changing (i.e. "converged")
 - if there have been large changes in value, extend the search to allow it to quiesce
 - rather than enforcing a strict depth limit, can be non-uniform
 - use a dynamic IS-CUTOFF(s) test
 - horizon effect
 - sometimes, enough dodging moves can be made to forestall a bad outcome so it occurs just beyond the depth limit (like moving a bishop back and forth to delay capture, or repeatedly checking the opponent's king)
 - delaying the inevitable it might change our decision if we knew this
 - hard to detect and mitigate

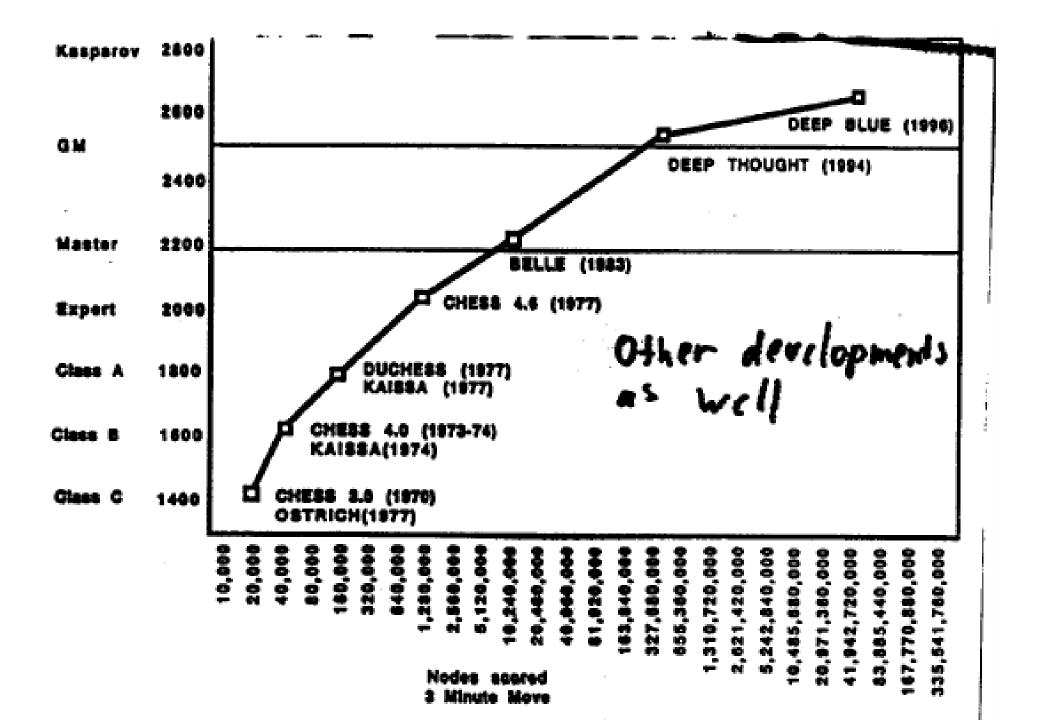
DeepBlue

- developed by IBM
- achieved grandmaster rating in 1990's
- defeated Gary Kasparov in 1997



- a supercomputer with **custom ASICs** for very fast α/β -Minimax search
 - 30-node IBM RS/6000 SP computer; 120 MHz and 1GB per proc.
 - 16 "chess chips" on each node, for generating moves and computing a board evaluation function
 - explored ~100 million moves/s, down to 10-12 ply (though non-uniform)
- included an **end-game database** (for example, once there are only 5 pieces left, lookup optimal moves in a pre-computed table)
- What did we learn about Intelligence?





Connect4

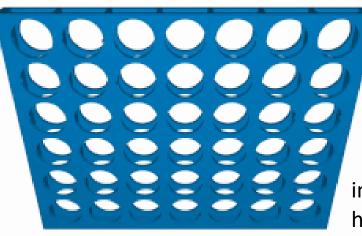


image obtained from
https://en.wikipedia.org/wiki/Connect_Four

- pieces are dropped in vertical columns; 4-in-a-row wins the game
 - here is an online app you can play around with: https://www.cbc.ca/kids/games/all/connect-4
- Challenge: <u>Can you come up with a board evaluation function for playing</u> <u>Connect4</u>?
 - it would not be hard to implement this on the command line (similar to tic-tac-toe
 - the State Space is much larger, so you would have to use a depth cutoff in the Minimax search and apply a board evaluation function to incomplete states
 - (try pausing the animation above and estimating the value of the state)

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a famous backgammon program called TDgammon (by Gary Tesauro) used Reinforcement Learning

CHANCE

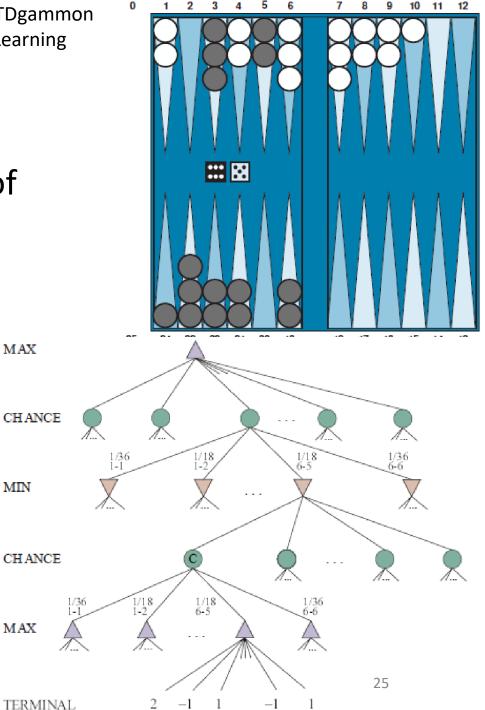
MAX

Expectiminimax

- stochastic games games with an element of chance (e.g. dice, cards...)
 - examples: backgammon, yahtze...
- can we apply minimax search?
 - MAX • yes, if we interleave min and max nodes with a level of chance nodes
 - CHANCE • at chance nodes, the score is the weighted sum over the children, weighted by probability, i.e. MIN "expected outcome"

Expectiminimax(s) = $u_1(s)$ if is a terminal node max{*Expectiminimax*(s')|s'∈succ(s)} if max node min{*Expectiminimax*(s')|s'∈succ(s)} if min node $\Sigma_{s' \in succ(s)} P(s') \cdot Expectiminimax(s')$ if chance node

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(Sec 5.4)

- instead of exhaustively exploring search tree, sample random paths ("rollouts") all the way to terminal states (end-games with defined utility)
- the value of a state is taken as the statistical *average* outcome of trajectories passing through it ("back-propagate" outcomes)
- also keep track of *n* (# trial trajectories passing through each node) and variance (σ²) at each state to assess certainty

function MONTE-CARLO-TREE-SEARCH(state) returns an action tree ← NODE(state) while IS-TIME-REMAINING() do leaf ← SELECT(tree) child ← EXPAND(leaf) result ← SIMULATE(child) BACK-PROPAGATE(result, child) return the move in ACTIONS(state) whose node has highest number of playouts

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- think of MCTS as an alternative to manually creating a board evaluation function
- estimate quality of each state (prob of winning) by simulating random game trajectories (playouts)

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6/6

16/53

10:18

2/11

3/4

0.3

- at each node, keep track of how many times it led to a win; more trajectories provide higher confidence
- can use these values to select children in minimax search
- select a node (game state) whose value is uncertain
- run simulation: play game to see outcome from that state 3/26
- back-propagation: update nodes along path with outcome 2/19/2023

- selection policy which states could use more sampling?
 - expansion vs. exploration
 - is it better to refine value estimate at good nodes, or increase certainty of bad nodes?
 - allow occasional sub-optimal choices for the sake of seeing how they turn out
- playout policy
 - there are many choices about how to make moves during simulation 60.79

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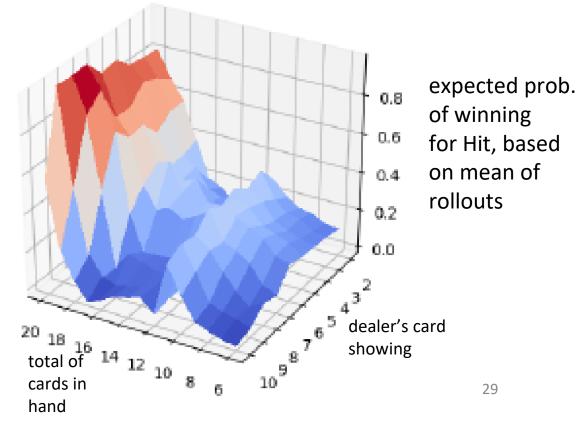
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- just making subsequent random moves is not realistic
- it helps to define an initial strategy to play against, even if weak

- using MCTS to learn strategy for Blackjack
 - simulate >10,000 random games to learn policy





AlphaGO

- GO is played with b/w stones on a 19x19 board
 - search space much larger than chess (bran. fact. starts at 361)
- from Google DeepMind, 2017
- after decades of attempts by other AI programs, AlphaGO finally beat the human GO world champion
- learns from *self-play* (bootstrapping), >100,000 games
- trains a *deep neural network* (14 conv. layers) to represent a value function (reinforcement learning, MCTS)
- reached grandmaster rating after 21 days (176 GPUs)



image from https://en.wikipedia.org/wiki/Go_(game)